Implicit Finite-Horizon Approximation and Efficient Optimal Algorithms for Stochastic Shortest Path

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Motivation

Many MDP models have been studied:

- Infinite horizon average reward model (Bartlett & Tewari, 2009; Jaksch et al., 2010)
- Infinite horizon discounted model (Even-Dar et al., 2003; Strehl et al., 2006)
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However, there are many real-world applications not modelled well by the above:

- Games (such as Go)
- Car navigation
- Robotic manipulation







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For these, Stochastic Shortest Path (SSP) is a better model.

- Episodic MDP with a goal state.
- Ends interaction only when the goal state is reached

Related Works

S: #states, A: #actions, K: #episodes, D: SSP-diameter c_{\min} : minimum cost, B_{\star} : maximum expected cost of optimal policy over all states T_{\star} : maximum expected hitting time of optimal policy starting from any state

- UC-SSP (Tarbouriech et al., 2020): $\tilde{\mathcal{O}}\left(DS\sqrt{\frac{D}{c_{\mathsf{min}}}AK} + S^2AD^2\right)$
- Bernstein-SSP (Cohen et al., 2020): $\tilde{\mathcal{O}}\left(B_{\star}S\sqrt{AK}+\sqrt{\frac{B_{\star}^{3}S^{2}A^{2}}{c_{\min}}}\right)$
- ULCVI (Cohen et al., 2021): $\tilde{\mathcal{O}}\left(B_{\star}\sqrt{SAK}+T_{\star}^{4}S^{2}A\right)$
- EB-SSP (Tarbouriech et al., 2021): $\tilde{\mathcal{O}}\left(B_{\star}\sqrt{SAK}+B_{\star}S^{2}A\right)$
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Techniques applied in previous works are quite different from each other, and some of these algorithms are fairly complicated.

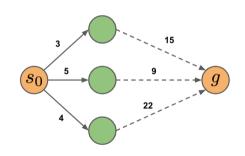
Our Results

Our contribution: A generic template for regret minimization algorithms in SSP. Using this template, we develop two algorithms:

S: #states, A: #actions, D: SSP-diameter, K: #episodes T_{\star} : expected hitting time of optimal policy, c_{\min} : minimum cost

	SVI-SSP	LCB-Advantage-SSP
Regret	$\int \tilde{\mathcal{O}}\left(B_{\star}\sqrt{SAK}+B_{\star}S^{2}A\right)$	$\tilde{\mathcal{O}}\left(B_{\star}\sqrt{SAK}+B_{\star}^{5}S^{2}A/c_{min}^{4}\right)$
Algorithm type	Model-based	Model-free (the first)

SSP Model: MDP $M = (S, A, s_{init}, g, c, P)$, only P is unknown.



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Learning Protocol

for $k = 1, \ldots, K$ do

learner starts in state $s_1^k = s_{\text{init}}, i \leftarrow 1$

while $s_i^k \neq g$ do

learner chooses action $a_i^k \in \mathcal{A}$ and observes states

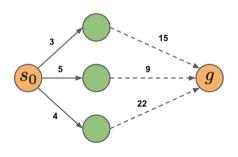
$$s_{i+1}^k \sim P(\cdot|s_i^k, a_i^k)$$

 $i \leftarrow i+1$

end

learner suffers cost $\sum_{i=1}^{l_k} c(s_i^k, a_i^k)$

end



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Notations:

- Policy π : maps state $s \in \mathcal{S}$ to an action $a \in \mathcal{A}$
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Objective: minimize regret w.r.t. the best proper policy in hindsight

$$R_{K} = \sum_{k=1}^{K} \left(\sum_{i=1}^{I_{k}} c(s_{k}^{i}, a_{k}^{i}) - V^{\pi^{\star}}(s_{0}) \right),$$

where $\pi^{\star} = \operatorname{argmin}_{\pi \in \Pi_{\operatorname{proper}}} V^{\pi}(s_0)$.

Generic Template

A General Algorithmic Template for SSP

```
Initialize: t \leftarrow 0, s_1 \leftarrow s_{\text{init}}, Q(s, a) \leftarrow c(s, a) for all (s, a). for k = 1, ..., K do
```

repeat

Increment time step $t \stackrel{+}{\leftarrow} 1$.

Take action $a_t = \operatorname{argmin}_a Q(s_t, a)$, suffer cost $c(s_t, a_t)$, and transit to s'_t .

Update Q (so that it satisfies Property 1 and Property 2).

if $s'_t \neq g$ then $s_{t+1} \leftarrow s'_t$; else $s_{t+1} \leftarrow s_{\mathsf{init}}$, break.

end

Record $T \leftarrow t$ (that is, the total number of steps).

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The key of analysis: Bounding the estimation error $Q^*(s_t, a_t) - Q(s_t, a_t)$. Issue: Relatively straightforward in a discounted setting or a finite-horizon setting, but becomes highly non-trivial in SSP.

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- We only need the optimal value functions of \widetilde{M} in the analysis:

$$Q_h^*(s,a) = c(s,a) + P_{s,a}V_{h-1}^*, \qquad V_h^*(s) = \min_{a} Q_h^*(s,a),$$

with $Q_0^{\star}(s,a) = 0$ for all (s,a).

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Lemma

For any value of H, $Q_H^\star(s,a) \leq Q^\star(s,a)$ holds for all (s,a). For any $\delta \in (0,1)$, if $H \geq \frac{4B_\star}{C_{\min}} \ln(2/\delta) + 1$, then $Q^\star(s,a) \leq Q_H^\star(s,a) + B_\star \delta$ holds for all (s,a).

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Similar approximation has been done explicitly before (Chen et al., 2021a; Chen et al., 2021b; Cohen et al., 2021)

To perform approximation implicitly, we need the following two properties of estimate Q (let Q_t be the value of Q at the beginning of time step t):

• **Property 1** (Optimism): with high probability, $Q_t(s, a) \leq Q^*(s, a)$ for all $(s, a), t \geq 1$.

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- **Property 1** (Optimism): with high probability, $Q_t(s, a) \leq Q^*(s, a)$ for all (s, a), $t \geq 1$.
- **Property 2** (Recursion): There exists a "bonus overhead" $\xi_H > 0$ and an absolute constant d > 0 such that the following holds with high probability:

$$egin{split} \sum_{t=1}^T (Q_h^\star(s_t, a_t) - Q_t(s_t, a_t)) &\leq \xi_H + \left(1 + rac{d}{H}
ight) \sum_{t=1}^T (V_{h-1}^\star(s_t) - Q_t(s_t, a_t))_+, \ \sum_{t=1}^T (Q^\star(s_t, a_t) - Q_t(s_t, a_t)) &\leq \xi_H + \left(1 + rac{d}{H}
ight) \sum_{t=1}^T (V^\star(s_t) - Q_t(s_t, a_t))_+, \end{split}$$

where $(x)_{+} = \max\{x, 0\}.$

Theorem

For any
$$\delta \in (0,1)$$
, if $H \ge \frac{4B_*}{c_{\min}} \ln(2/\delta) + 1$, then the template ensures (with high probability) $R_K = \tilde{\mathcal{O}}\left(\sqrt{B_*\mathcal{C}_K} + B_* + \delta\mathcal{C}_K + \xi_H\right)$, where $\mathcal{C}_K = \sum_{k=1}^K \sum_{i=1}^{I_k} c(s_i^k, a_i^k)$.

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Now if we ensure $\xi_H = \tilde{\mathcal{O}}\left(\sqrt{B_{\star}SAC_{K}}\right)$ (with appropriate bonus), then $R_{K} = \tilde{\mathcal{O}}\left(B_{\star}\sqrt{SAK}\right)$.

No explicit implementation of \widetilde{M} is required!

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Update Q(s, a) for logarithmically many times for each (s, a). When updating Q(s, a), we apply the following update rule:

$$Q(s,a) \leftarrow \max\left\{c(s,a) + \bar{P}_{s,a}V - b, Q(s,a)\right\},$$

where \bar{P} is the empirical transition, $b \approx \max\left\{7\sqrt{\frac{\mathbb{V}(\bar{P}_{s,a},V)}{n}},\frac{49B_{\star}}{n}\right\}$ (Zhang et al., 2021).

Theorem

SVI-SSP satisfies Property 1 and Property 2 with d=1 and $\xi_H = \tilde{\mathcal{O}} \left(\sqrt{B_{\star} SAC_K} + B_{\star} S^2 A + \delta C_K \right)$.

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- Lower time complexity of updates: SVI-SSP: $\tilde{\mathcal{O}}(B_{\star}S^2A/c_{\min})$, EB-SSP: $\tilde{\mathcal{O}}(B_{\star}^2S^5A/c_{\min}^2)$, ULCVI: $\tilde{\mathcal{O}}(S^2AT_{\star}K)$

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Our implicit finite horizon analysis is the key to achieve sparse updates.

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Update Q(s,a) for logarithmically many times for each (s,a). Inspired by (Zhang et al., 2020), we update Q(s,a) with the following variance reduced update rule (approximately)

$$Q(s,a) \leftarrow \max \left\{ c(s,a) + \frac{1}{n} \sum_{i=1}^n V^{\mathsf{ref}}(s'_{t_i}) + \frac{1}{m} \sum_{i=1}^m \left(V(s'_{t'_i}) - V^{\mathsf{ref}}(s'_{t'_i}) \right) - b, Q(s,a) \right\},$$

where m is the number of samples in current stage, and n is the number of samples up to current stage, and $V(s) = \min_a Q(s, a)$.

Theorem

 $\operatorname{LCB-Advantage-SSP}$ satisfies Property 1 and Property 2 with d=3 and

$$\xi_H = \tilde{\mathcal{O}}\left(\sqrt{B_{\star}SAC_K} + \frac{B_{\star}^2H^3S^2A}{c_{\min}}\right).$$

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LCB-Advantage-SSP ensures
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• To make it parameter-free, we try logarithmically many different values of parameters simultaneously, each leading to a different update rule for Q and V^{ref} .

Summary

Our contribution: A generic template for regret minimization algorithms in SSP. Using this template, we develop two algorithms:

S: #states, A: #actions, D: SSP-diameter, K: #episodes T_{\star} : expected hitting time of optimal policy, c_{\min} : minimum cost

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Algorithm type	Model-based	Model-free (the first)